

§1: Functor of points:

Not<sup>n</sup>:  $\mathcal{C}$  category  $\mathcal{C}^{opp}$  opposite category  
 same objects, arrows reversed  
 $Hom_{\mathcal{C}^{opp}}(x, y) = Hom_{\mathcal{C}}(y, x)$   
 $Hom_{\mathcal{C}^{opp}} = Hom_{\mathcal{C}}$

Def<sup>n</sup>: Let  $\mathcal{C}, \mathcal{S}$  be two categories.  
 A contravariant functor  $F$  from  $\mathcal{C}$  to  $\mathcal{S}$  is a functor  
 $F: \mathcal{C}^{opp} \rightarrow \mathcal{S}$   
 $F: Ob(\mathcal{C}) \rightarrow Ob(\mathcal{S})$   
 and for every  $\varphi: x \rightarrow y \mapsto F(\varphi): F(y) \rightarrow F(x)$   
 (composition is reversed as well)

Def<sup>n</sup>: Let  $\mathcal{C}$  be a category.  
 (1) A presheaf of "set" on  $\mathcal{C}$  is a contravariant functor  $F$  from  $\mathcal{C}$  to  $Set$  (Set of sets).  
 (2) The category of presheaves is  $PSh(\mathcal{C})$ .

Example: Functor of points For  $U \in Ob(\mathcal{C})$

$$h_U: \mathcal{C} \rightarrow Set$$

contravariant  $X \mapsto Hom_{\mathcal{C}}(X, U)$

also denoted  $h_U: \mathcal{C}^{opp} \rightarrow Set$

Also called representable presheaf associated to  $U$ .

Note that  $\phi: U \rightarrow V$  in  $\mathcal{C}$  we have  $h(\phi): h_U \rightarrow h_V$   
 morphism (defined by  $h_U(T) \rightarrow h_V(T)$   
 $Hom(T, U) \xrightarrow{h(\phi)} Hom(T, V)$   
 $T \rightarrow U \mapsto T \rightarrow V$ )

This gives us a functor to the category of Functors

$$h: \mathcal{C} \rightarrow Fun(\mathcal{C}^{opp}, Set) = PSh(\mathcal{C})$$

"big category" ← set theoretic remark...

Scheme (Sch)  
 $\mathbb{Z}$   
 $\mathcal{C} = (Sch/S) \vee \mathcal{C} \in (Sch/S)$   
 $h_U: (Sch/S)^{opp} \rightarrow Set$   
 Functor of points of scheme  $U$ .  
Ex: Show that all  $k$ -points of  $S$  scheme  $U$  are  $h_U(k)$ .

(Yoneda lemma): Let  $U, V \in Ob(\mathcal{C})$ . Given any morphism of functors  $s: h_U \rightarrow h_V$  there is a unique morphism  $\phi: U \rightarrow V$  such that  $h(\phi) = s$ .

In other words, the functor  $h$  is fully faithful.

More generally, given any contravariant functor  $F$  and any object  $U$  of  $\mathcal{C}$  we have a natural bijection

$$Hom_{PSh(\mathcal{C})}(h_U, F) \rightarrow F(U)$$

$$s \mapsto s_U(id_U)$$

PF:  $s_U: h_U(U) \rightarrow h_V(U)$  Take  $\phi = s_U(id_U)$   
 $Hom(U, U) \xrightarrow{id} Hom(U, V)$   
 $id \mapsto s_U(id_U)$

For the second statement... Given  $\Sigma \in F(U)$ , define  $s$  by  $s_V: h_V(V) \rightarrow F(V)$   
 $f: V \rightarrow U \mapsto F(f)(\Sigma)$

Def<sup>n</sup>: A contravariant functor

□

Def<sup>n</sup>: A contravariant functor  $F: \mathcal{C} \rightarrow \text{Sets}$  is said to be representable if it is isomorphic to the functor of points  $h_U$  for some  $U$  of  $\mathcal{C}$ .  $\square$

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$$F \cong h_U$$

Remarks:  $\mathcal{C}: F: \mathcal{C}^{\text{opp}} \rightarrow \text{Sets}$  is a rep<sup>n</sup> functor.

Choose an object  $U$  of  $\mathcal{C}$  and an iso  $s: h_U \rightarrow F$ . Then Yoneda lemma guarantees that the pair  $(U, s)$  is unique upto a unique morphism.  $U$  is called object representing  $F$ .

By Yoneda lemma the transformation  $s$  corresponds to a unique element  $\xi \in F(U)$

↑  
universal object.

It has the property that for  $V \in \text{Ob}(\mathcal{C})$ , the map  $\text{Mor}_{\mathcal{C}}(V, U) \rightarrow F(V)$

$(f: V \rightarrow U) \mapsto F(f)(\xi)$  is a bijection.

Thus  $\xi$  is universal in the sense that every element of  $F(V)$  is equal to the image of  $\xi$  via  $F(f)$ , for a unique morphism  $f: V \rightarrow U$  in  $\mathcal{C}$ .

let  $\mathcal{C} = (\text{Sch}) / (\text{Sch}/s)$ .

Def<sup>n</sup>: A presheaf  $F: (\text{Sch}) \rightarrow \text{Sets}$  is said to be a sheaf in the Zariski topology if for every scheme  $T$  and every open covering  $T = \cup_{i \in I} U_i$  and for any collection of elements  $\xi_i \in F(U_i)$  s.t.  $\xi_i|_{U_i \cap U_j} = \xi_j|_{U_i \cap U_j}$  there

open covering  $\{U_i, i \in I\}$  of elements  $\xi_i \in F(U_i)$  s.t.  $\xi_i|_{U_i \cap U_j} = \xi_j|_{U_i \cap U_j}$  there exist a unique element  $\xi \in F(T)$  s.t.  $\xi|_{U_i} = \xi_i|_{U_i}$  in  $F(U_i)$ .

Exercise: Show that  $h_U$  for any scheme  $U$  is a sheaf in the Zariski topology. In particular, deduce that for a presheaf to be representable it must be a sheaf in the Zariski topology.

Moduli problem:

Aim: We want to classify a collection of objects (in our case a collection of curves) over an algebraically closed field upto certain equivalence (isomorphism)

$M_g :=$  set of curves of genus  $g$  /  $k$  upto isomorphism.

structure of a scheme over this set.

- Notion of family.

Geometry of  $M_g$  to relate isomorphism classes of curves in some way.

"limits"

$X$  ("deformation")

Remark: smooth projective curves /  $\mathbb{C}$

Thm (Torelli):  $X, Y$  are curves as above. Then

$X \cong Y$  iff  $Jac(X) \cong Jac(Y)$   
 (isomorphism in category of schemes)

structure of  $Pic^0(X) :=$  set of degree zero line bundles on  $X$ .  
 an abelian variety.

$M_g :=$  Notion of a family

Family:  $\pi: X \rightarrow B$  morphism of schemes  
 s.t.  $\pi^{-1}(b)$  is a scheme.  $b \in B$  point

This is too general (useless), fibers have nothing in common  
 $\pi: X \rightarrow B$  be a family  $b \in B$  be a closed point.

New family  $\pi': (X \setminus \pi^{-1}(b)) \sqcup Y \rightarrow B$

$$\begin{array}{ccc} Y & \hookrightarrow & b \\ X \setminus \pi^{-1}(b) & \hookrightarrow & B \setminus b \end{array} \quad \text{Y any scheme}$$

We have to put some conditions on the type of morphisms to be allowed.

Ans: The best current candidate is flatness.

- local triviality of family  $X \times B \rightarrow B$  is bad in the locally coarse Zariski topology
- local completions  $\xleftrightarrow{\text{local}} \text{analytic rings}$   
 $\xleftrightarrow{\text{GAL, A}}$  only smooth family...

To pose a moduli questions:

1. A class of geometric objects (scheme, line bundles, vector bundles, perfect complexes)
2. A notion of family.
3. A notion of equivalence

Def<sup>n</sup>: A moduli problem for a class  $\mathcal{P}$  of objects in a category  $\mathcal{C}$  is a contravariant functor

$$F_{\mathcal{P}}: \mathcal{C} \rightarrow \text{sets}$$

that associates to each  $B \in \mathcal{C}$  the set of isomorphism classes of families of  $\mathcal{P}$ -objects.

To each morphism  $f: B' \rightarrow B$  it associates the set sending a family  $X \rightarrow B$  to the pullback family  $f^*(X) \rightarrow B'$

Def: A object  $M_{\mathcal{P}} \in \mathcal{C}$  that represents the functor  $F_{\mathcal{P}}$  is called a fine moduli space.

Def: A object  $M_p \in \mathcal{C}$  over represents the point  $p$ , is called a fine moduli space.

Rmk. Yoneda  $\Rightarrow$  Fine moduli space is unique upto unique isomorphism.

Back to our notion of family in case of schemes.

- At least dimension of fibers should be same
- certain (numerical) invariants be constant in the fibers of the family (genus, Hilbert polynomial)

Prop<sup>n</sup>: Let  $f: X \rightarrow Y$  be a flat morphism of schemes of finite type over a field  $k$ . For any point  $x \in X$ , let  $y = f(x)$ . Then  $\dim_x(X_y) = \dim_x X - \dim_y Y$ .

Proof: Reduction to the case  $y$  is a closed point of  $Y$ .  
 $\dim_x X := \dim \mathcal{O}_{X,x}$

Take  $Y' = \text{Spec } \mathcal{O}_{Y,y}$ .  $x \in X' \rightarrow X_x$   
 Base change  $\text{flat} \downarrow \quad \downarrow \text{flat}$

Check:  $\dim_x X' = \dim_x X$   
 $\dim_x |X'_y| = \dim_x(X_y)$  (closed point)  
 $\dim_y Y = \dim_y Y'$

Now we use induction on  $\dim Y$ .

If  $\dim Y = 0$ . (Ex.)  $X_y$  is defined by a nilpotent ideal in  $X$ .  
 we have  $\dim_x(X_y) = \dim_x X$  and  $\dim_y Y = 0$ .

If  $\dim Y > 0$ , base change to  $Y_{\text{red}}$  (dim don't change)  
 assume  $Y$  is reduced.

Then find an element  $t \in \mathfrak{m}_y \in \mathcal{O}_{Y,y}$  s.t  $t$  is a non-zero divisor

let  $Y' = \text{Spec } \mathcal{O}_{Y,y}/(t)$  and make the base change to  $Y' \rightarrow Y$ .

(Ex)  $\dim Y' = \dim Y - 1$ ,

Since  $f$  is flat,  $f^\# t \in \mathfrak{m}_x$  is not a zero divisor  
 so for the same reason  $\dim_x X' = \dim_x X - 1$ .

But  $X_y$  does not change under base change to  $Y'$   $\square$ .

Corollary: let  $f: X \rightarrow Y$  be a flat morphism of schemes of  $f +$  over a field  $k$  and assume that  $Y$  is irreducible.

TFAE:

- (i) every irreducible component of  $X$  has  $\dim$  equal to  $\dim Y + n$ .
- (ii) for any point  $y \in Y$  (closed or not), every irreducible component of the fibre  $X_y$  has dimension  $n$ .

Pf: (i)  $\Rightarrow$  (ii) Given  $y \in Y$ ,  $Z \subseteq X_y$  irred comp. and let  $x \in Z$  be a closed pt. st.  $x \notin Z'$  for any other irred component of  $X_y$ .

Applying the previous result

$$\dim_x Z = \dim_x X - \dim_y Y.$$

$$\text{Now } \dim_x Z = \dim Z \quad (x \text{ is closed})$$

Since  $Y$  is irreducible,  $X$  is equi-dimensional + fl.

$$\dim_x X = \dim X - \dim \overline{\{x\}} \quad [E.x.]$$

$$\dim_y Y = \dim Y - \dim \overline{\{y\}}$$

$x$  was closed in  $X_y$ ,  $k(x)$  is a finite algebraic ext<sup>n</sup> of  $k(y)$  and so

$$\dim \overline{\{x\}} = \dim \overline{\{y\}}$$

Combining this with (i) we get  $\dim Z = n$ .

(ii)  $\Rightarrow$  (i) Exercise -

Prop<sup>n</sup>: let  $f: X \rightarrow Y$  be a morphism of schemes, with  $Y$  integral and regular of  $\dim 1$ . Then  $f$  is flat iff every associated point  $x \in X$ , maps to the generic point of  $Y$ . In particular, if  $X$  is reduced,  $f$  is flat iff every irreducible component of  $X$  dominates  $Y$ .